

Wikipedia in the Pocket

Indexing Technology for Near-duplicate Detection and High Similarity Search

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ABSTRACT

We develop and implement a new indexing technology which allows us to use complete (and possibly very large) documents as queries, while having a retrieval performance comparable to a standard term query. Our approach aims at retrieval tasks such as near-duplicate detection and high similarity search. To demonstrate the performance of our technology we have compiled the search index “Wikipedia in the Pocket”, which contains about 2 million English and German Wikipedia articles.¹ This index—along with a search interface—fits on a conventional CD (0.7 gigabyte). The ingredients of our indexing technology are similarity hashing and minimal perfect hashing.

Categories and Subject Descriptors: E.2 [Data]: Data Storage Representations—*Hash-table representations*; H.3.1 [Information Storage and Retrieval]: Content Analysis and Indexing—*Abstracting methods, Indexing methods*

General Terms: Algorithms, Performance

Keywords: hash-based indexing, fuzzy-fingerprinting, near-duplicate detection

Similarity Hashing

We use tailored similarity hash functions $h_\varphi : D \rightarrow U, U \subset \mathbb{N}$, to map documents with high pairwise similarity onto the same integer hashcode with high probability. Fuzzy-fingerprinting, which was specifically designed for text-based information retrieval, serves as a means for the construction of such hash functions [2]. It is based on the definition of a small number of $k, k \in [10, 100]$, prefix equivalence classes. A prefix class, for short, contains all terms starting with the same prefix. The computation of $h_\varphi(d)$ happens in the following steps: (1) Computation of \mathbf{pf} , a k -dimensional vector that quantifies the distribution of the index terms in d with respect to the prefix classes. (2) Normalization of \mathbf{pf} using a corpus that provides a representative cross-section of the source language, and computation of $\Delta_{\mathbf{pf}} = (\delta_1, \dots, \delta_k)^T$, the vector of deviations to the expected distribution.² (3) Fuzzification of $\Delta_{\mathbf{pf}}$ by projecting the exact deviations according to diverse fuzzification schemes.

Index Data Structure

Given a similarity hash function $h_\varphi : D \rightarrow U$, an index can be directly constructed by means of a hash table \mathcal{T} along with a standard

¹A Wikipedia snapshot from November 4th, 2006 was indexed.

²The British National Corpus is used as reference.

hash function $h : U \rightarrow \{1, \dots, |\mathcal{T}|\}$; h maps the universe of hashcodes, U , onto the $|\mathcal{T}|$ storage positions. To index a document d a reference to it is added to the bucket at storage position $h(h_\varphi(d))$; likewise all documents $D' \subseteq D$ similar to d can be retrieved from this bucket.

In the case of Wikipedia in the Pocket all documents to be indexed are given and hence we can construct a minimal perfect hash function h . If h is perfect no two hashcodes will be mapped to the same storage position in \mathcal{T} , moreover, if h is minimal no additional storage positions are needed to ensure the former. Hence the resulting index is space optimal and allows the retrieval of similar documents in $O(1)$ time.

However, space optimal hashing entails also shortcomings: The minimal perfect hash function h has a $4.6 \cdot |D|$ byte memory overhead to represent the perfect mapping [1]. Furthermore, keys unknown prior to the construction of h will also map to a storage position in \mathcal{T} . Therefore the preimage of h must be stored twice, in h , as well as at the respective storage positions of \mathcal{T} . The former allows a perfect mapping, the latter prevents false matches. We have relaxed the latter and store small checksums: The probability of a false match between two hashcodes $x, y \in U, x \neq y$, that map to the same storage position is $p = (\frac{1}{2})^c$ where c denotes the bit length of the checksum—in our case $c = 16$ bit.

Figure 1 shows benchmark data of the index. The table lists the sizes of the English and German index, the diagram shows precision and recall curves dependent on similarity thresholds.



Figure 1: Benchmark data of Wikipedia in the Pocket

References

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